

Monoidal

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Monoid

Definition (Monoid)

Let C be a set with $1 \in C$ and $\cdot : C \times C \rightarrow C$, such that

$$(a \cdot b) \cdot c = a \cdot (b \cdot c)$$

$$a \cdot 1 = 1 \cdot a = a$$

then $\langle C, \cdot, 1 \rangle$ is a *monoid*.

Examples

- Strings with concatenation and the empty string

$(\text{"foo"} \text{++} \text{"bar"}) \text{++} \text{"baz"} = \text{"foo"} \text{++} (\text{"bar"} \text{++} \text{"baz"})$

$\text{"foo"} \text{++} \text{""} = \text{"foo"} = \text{""} \text{++} \text{"foo"}$

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$$\text{"foo"} \text{ ++ } \text{""} = \text{"foo"} = \text{""} \text{ ++ } \text{"foo"}$$

- Integers with addition and 0

$$(123 + 456) + 789 = 123 + (456 + 789)$$
$$42 + 0 = 42 = 0 + 42$$

Examples

- Strings with concatenation and the empty string

$$("foo" ++ "bar") ++ "baz" = "foo" ++ ("bar" ++ "baz")$$

$$"foo" ++ "" = "foo" = "" ++ "foo"$$

- Integers with addition and 0

$$(123 + 456) + 789 = 123 + (456 + 789)$$

$$42 + 0 = 42 = 0 + 42$$

- Sets with union and the empty set

$$(\{A, B\} \cup \{B, C\}) \cup \{A, D\} = \{A, B\} \cup (\{B, C\} \cup \{A, D\})$$

$$\{A\} \cup \emptyset = \{A\} = \emptyset \cup \{A\}$$

Strict Monoidal Categories

Definition

Let C be a category with $1 \in C$ and a bifunctor $\otimes : C \times C \rightarrow C$, such that

$$(X \otimes Y) \otimes Z = X \otimes (Y \otimes Z)$$

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and

$$a : (X \otimes Y) \otimes Z \xrightarrow{\sim} X \otimes (Y \otimes Z) = id$$

$$\iota : 1 \otimes 1 \xrightarrow{\sim} 1 = id$$

then $\langle C, \otimes, a, 1, \iota \rangle$ is a *strict monoidal category*.

Example

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Let C be a category, then $End(C)$ is a strict monoidal category.

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Proof.

- $f \otimes g = f \circ g$
- $1 = id_C$

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- $\forall X, Y, Z \in C. (X \otimes Y) \otimes Z = X \otimes (Y \otimes Z)$
by composition laws.

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- $1 = id_C$
- $\forall X, Y, Z \in C. (X \otimes Y) \otimes Z = X \otimes (Y \otimes Z)$
by composition laws.
- $1 \otimes 1 = 1$ by identity laws.



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Let \mathbb{F} be a field, then $Mat(\mathbb{F})$ is a strict monoidal category

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- $n \otimes m = n \cdot m$, $f \otimes g$ by Kronecker product
- 1 by single 1-matrix

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Proof.

- $n \otimes m = n \cdot m$, $f \otimes g$ by Kronecker product
- 1 by single 1-matrix
- $\forall X, Y, Z \in C. (X \otimes Y) \otimes Z = X \otimes (Y \otimes Z)$
by properties of Kronecker product
- $1 \otimes 1 = 1$ obvious



Non-Strict Monoidal Categories

Definition (Monoidal Category)

Let \mathcal{C} be a category with $1 \in \mathcal{C}$, a bifunctor $\otimes : \mathcal{C} \times \mathcal{C} \rightarrow \mathcal{C}$, and

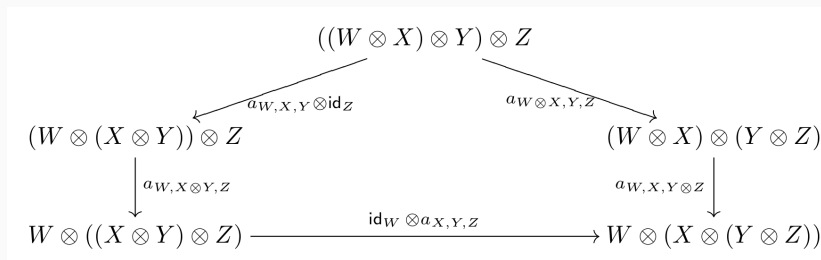
$$a_{X,Y,Z} : (X \otimes Y) \otimes Z \xrightarrow{\sim} X \otimes (Y \otimes Z)$$

$$\iota : 1 \otimes 1 \xrightarrow{\sim} 1$$

satisfying the *pentagon* and the *unit axiom*, then $\langle \mathcal{C}, \otimes, a, 1, \iota \rangle$ is a *monoidal category*

Pentagon Axiom

The diagram



commutes

The functors

$$L_1 : X \mapsto 1 \otimes X$$

$$R_1 : X \mapsto X \otimes 1$$

are autoequivalences of \mathcal{C} .

Unit constraints

Define natural isomorphisms

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$$r_X : X \otimes 1 \rightarrow X$$

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$$r_X : X \otimes 1 \rightarrow X$$

such that

$$L_1(l_X) = 1 \otimes (1 \otimes X) \xrightarrow{a_{1,1,X}^{-1}} (1 \otimes 1) \otimes X \xrightarrow{\iota \otimes \text{id}_X} 1 \otimes X$$

$$R_1(r_X) = (X \otimes 1) \otimes 1 \xrightarrow{a_{X,1,1}} X \otimes (1 \otimes 1) \xrightarrow{\text{id}_X \otimes \iota} X \otimes 1$$

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$$R_1(r_X) = (X \otimes 1) \otimes 1 \xrightarrow{a_{X,1,1}} X \otimes (1 \otimes 1) \xrightarrow{\text{id}_X \otimes \iota} X \otimes 1$$

then l_X is the left and r_X is the right *unit constraint*

Unit constraints

For all $X \in \mathcal{C}$

$$l_{1 \otimes X} = id_1 \otimes l_X$$

$$r_{X \otimes 1} = r_X \otimes id_1$$

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$$l_{1 \otimes X} = id_1 \otimes l_X$$

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Proof.

$$\begin{array}{ccc} \mathbf{1} \otimes (\mathbf{1} \otimes X) & \xrightarrow{\mathbf{1} \otimes l_X} & \mathbf{1} \otimes X \\ l_{\mathbf{1} \otimes X} \downarrow & & \downarrow l_X \\ \mathbf{1} \otimes X & \xrightarrow{l_X} & X. \end{array}$$

□

Triangle

The diagram

$$\begin{array}{ccc} (X \otimes \mathbf{1}) \otimes Y & \xrightarrow{a_{X, \mathbf{1}, Y}} & X \otimes (\mathbf{1} \otimes Y) \\ & \searrow r_X \otimes \text{id}_Y & \swarrow \text{id}_X \otimes l_Y \\ & X \otimes Y & \end{array}$$

commutes

Triangle

Proof.

$$\begin{array}{ccc} ((X \otimes \mathbf{1}) \otimes \mathbf{1}) \otimes Y & \xrightarrow{a_{X,1,1} \otimes \text{id}_Y} & (X \otimes (\mathbf{1} \otimes \mathbf{1})) \otimes Y \\ \downarrow a_{X \otimes \mathbf{1}, \mathbf{1}, Y} & \swarrow r_X \otimes \text{id}_{\mathbf{1}} \otimes \text{id}_Y \quad \nwarrow (\text{id}_X \otimes \iota) \otimes \text{id}_Y & \downarrow a_{X, \mathbf{1} \otimes \mathbf{1}, Y} \\ & (X \otimes \mathbf{1}) \otimes Y & \\ & \downarrow a_{X, \mathbf{1}, Y} & \\ & X \otimes (\mathbf{1} \otimes Y) & \\ \downarrow a_{X \otimes \mathbf{1}, \mathbf{1} \otimes Y} & \swarrow r_X \otimes \text{id}_{\mathbf{1} \otimes Y} \quad \nwarrow \text{id}_X \otimes (\iota \otimes \text{id}_Y) & \downarrow a_{X, \mathbf{1} \otimes \mathbf{1}, Y} \\ (X \otimes \mathbf{1}) \otimes (\mathbf{1} \otimes Y) & \xrightarrow{\text{id}_X \otimes l_{\mathbf{1} \otimes Y}} & X \otimes ((\mathbf{1} \otimes \mathbf{1}) \otimes Y) \\ \downarrow a_{X, \mathbf{1}, \mathbf{1} \otimes Y} & \uparrow \text{id}_X \otimes l_{\mathbf{1} \otimes Y} & \downarrow \text{id}_X \otimes a_{\mathbf{1}, \mathbf{1}, Y} \\ & X \otimes (\mathbf{1} \otimes (\mathbf{1} \otimes Y)) & \end{array}$$

□

Triangles

The diagrams

$$\begin{array}{ccc} (\mathbf{1} \otimes X) \otimes Y & \xrightarrow{a_{\mathbf{1}, X, Y}} & \mathbf{1} \otimes (X \otimes Y) \\ & \searrow l_X \otimes \text{id}_Y & \swarrow l_{X \otimes Y} \\ & X \otimes Y & \end{array}$$

$$\begin{array}{ccc} (X \otimes Y) \otimes \mathbf{1} & \xrightarrow{a_{X, Y, \mathbf{1}}} & X \otimes (Y \otimes \mathbf{1}) \\ & \searrow r_{X \otimes Y} & \swarrow \text{id}_X \otimes r_Y \\ & X \otimes Y & \end{array}$$

commute.

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commute.

Proof.
similar



Proposition

Corollary

In any monoidal category $l_1 = r_1 = \iota$

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Proof.

Set $X = Y = 1$ in the second triangle. We have

$$l_1 \otimes id_1 = l_{1 \otimes 1} \circ a_{1,1,1} = (id_1 \otimes l_1) \circ a_{1,1,1}$$

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By unit constraint

$$(id_1 \otimes l_1) \circ a_{1,1,1} = \iota \otimes id_1$$

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Hence

$$r_1 \otimes id_1 = l_1 \otimes id_1 = \iota \otimes id_1$$

Unique unit

The unit object in a monoidal category is unique up to a unique isomorphism.

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Proof.

Let $\langle 1_1, \iota_1 \rangle$, $\langle 1_2, \iota_2 \rangle$ be two unit objects with unit constraints $\langle r_1, l_1 \rangle$, $\langle r_2, l_2 \rangle$.

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We have $\eta := l_1 (1_2) \circ (r_2 (1_1))^{-1} : 1_1 \xrightarrow{\sim} 1_2$

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$$\begin{array}{ccc} \mathbf{1} \otimes \mathbf{1} & \xrightarrow{b \otimes b} & \mathbf{1} \otimes \mathbf{1} \\ \downarrow \iota & & \downarrow \iota \\ \mathbf{1} & \xrightarrow{b} & \mathbf{1} \end{array}$$

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Note that

$$\begin{array}{ccc} \mathbf{1} \otimes \mathbf{1} & \xrightarrow{c \otimes \text{id}_1} & \mathbf{1} \otimes \mathbf{1} \\ \downarrow \iota & & \downarrow \iota \\ \mathbf{1} & \xrightarrow{c} & \mathbf{1} \end{array}$$

Example

Sets with the Cartesian product and the singleton set are a monoidal category.

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$\mathbb{K}\text{-Vec}$ with $\otimes = \otimes_{\mathbb{K}}$ and $1 = \mathbb{K}$ is a monoidal category.

Theorem (Mac Lane)

Any monoidal category is monoidally equivalent to a strict monoidal category

Mac Lane's strictness theorem

Theorem (Mac Lane)

Any monoidal category is monoidally equivalent to a strict monoidal category

Proof.

Tensor Categories - P. 37



Rigid Monoidal Categories

Definition (Left Dual)

Let C be a monoidal category and $X \in C$. If there exist

$ev_X : X^* \otimes X \rightarrow 1$ and $coev_X : 1 \rightarrow X \otimes X^*$ such that

$$\begin{aligned} X &\xrightarrow{coev_X \otimes id_X} (X \otimes X^*) \otimes X \xrightarrow{a_{X, X^*, X}} X \otimes (X^* \otimes X) \xrightarrow{id_X \otimes ev_X} X \\ X^* &\xrightarrow{id_{X^*} \otimes coev_X} X^* \otimes (X \otimes X^*) \xrightarrow{a_{X^*, X, X^*}^{-1}} (X^* \otimes X) \otimes X^* \xrightarrow{ev_X \otimes id_{X^*}} X^* \end{aligned}$$

then X^* is the *left dual* of X .

Definition (Right Dual)

Let \mathcal{C} be a monoidal category and $X \in \mathcal{C}$. If there exist $ev'_X : X \otimes *X \rightarrow 1$ and $coev'_X : 1 \rightarrow *X \otimes X$ such that

$$\begin{aligned} X &\xrightarrow{id_X \otimes coev'_X} X \otimes (*X \otimes X) \xrightarrow{a_{X, *X, X}^{-1}} (X \otimes *X) \otimes X \xrightarrow{ev'_X \otimes id_X} X \\ *X &\xrightarrow{coev'_X \otimes id_{*X}} (*X \otimes X) \otimes *X \xrightarrow{a_{*X, X, *X}} *X \otimes (X \otimes *X) \xrightarrow{id_{*X} \otimes ev'_X} X \end{aligned}$$

then $*X$ is the *right dual* of X .

Example

\mathbb{K} -Vec is rigid.

- Dual of V is the dual vector space V^*
- $ev_V : V^* \otimes V \rightarrow \mathbb{K}$ by contraction
- $coev_V : \mathbb{K} \rightarrow V^* \otimes V$ by embedding

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If $X \in \mathcal{C}$ has a left / right dual, then it is unique up to a unique isomorphism.

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If $X \in \mathcal{C}$ has a left / right dual, then it is unique up to a unique isomorphism. Define $\alpha : X_1^* \rightarrow X_2^*$ by

$$X_1^* \xrightarrow{id_{X_1^*} \otimes c_2} X_1^* \otimes (X \otimes X_2^*) \xrightarrow{a_{X_1^*, X, X_2^*}^{-1}} (X_1^* \otimes X) \otimes X_2^* \xrightarrow{e_1 \otimes id_{X_2^*}} X_2^*$$

and $\beta : X_2^* \rightarrow X_1^*$ similar.

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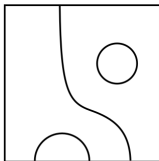
Show $\beta \circ \alpha$ and $\alpha \circ \beta$ are *id*.

$$\begin{array}{ccccc}
 X_1^* & \xrightarrow{id \otimes c_1} & X_1^* \otimes X \otimes X_1^* & & \\
 id \otimes c_2 \downarrow & & id \otimes c_2 \otimes id \downarrow & \searrow id & \\
 X_1^* \otimes X \otimes X_2^* & \xrightarrow{id \otimes c_1} & X_1^* \otimes X \otimes X_2^* \otimes X \otimes X_1^* & \xrightarrow{id \otimes e_2 \otimes id} & X_1^* \otimes X \otimes X_1^* \\
 e_1 \otimes id \downarrow & & e_1 \otimes id \downarrow & & \downarrow e_1 \otimes id \\
 X_2^* & \xrightarrow{id \otimes c_1} & X_2^* \otimes X \otimes X_1^* & \xrightarrow{e_2 \otimes id} & X_1^*
 \end{array}$$

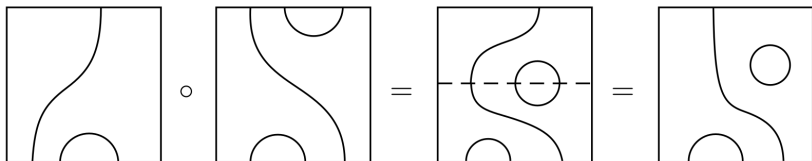
Temperley-Lieb categories

Definition (simple Temperley-Lieb diagram)

Let m, n be non-negative. Consider unit square with m and n points on top and bottom. A *simple Temperley-Lieb diagram* consists of smooth, non-crossing arcs between top and bottom.



Composition



The *Temperley-Lieb category* is a category where

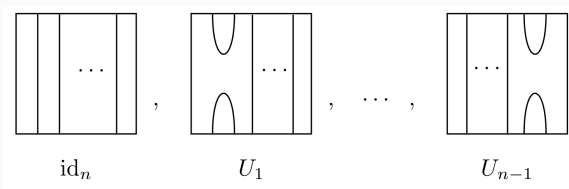
- Objects are non-negative \mathbb{N}
- $\text{Hom}(m, n)$ is the \mathbb{F} -linear span of diagrams from m to n , modulo d -equivalence.

The diagrammatic equation shows three square boxes connected by equals signs. The first box contains a vertical strand on the right, a semi-circle at the top, and two circles on the left. The second box contains the same strand and semi-circle, but with only one circle on the left. The third box contains the strand and semi-circle with no circles. The equation is labeled with d between the first and second boxes, and d^2 between the second and third boxes, indicating that each circle is worth a factor of d .

$$\boxed{\text{strand, semi-circle, 2 circles}} = d \boxed{\text{strand, semi-circle, 1 circle}} = d^2 \boxed{\text{strand, semi-circle}}$$

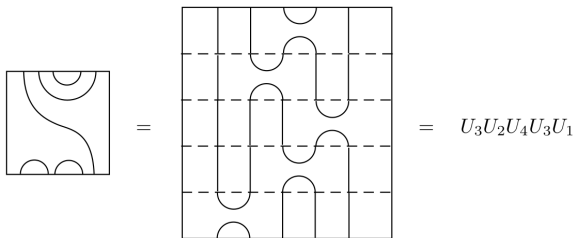
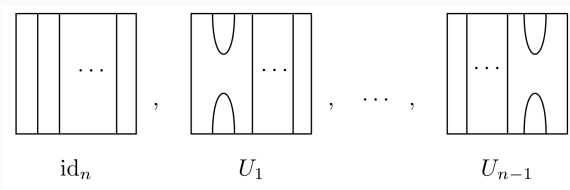
Generators

All morphisms $f \in \text{Hom}(m, n)$ can be written canonically



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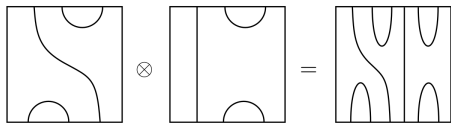
All morphisms $f \in \text{Hom}(m, n)$ can be written canonically



Strict Monoidal

The Temperley-Lieb category is a strict monoidal category where

- 1 is the 0-diagram
- for all objects $a, b \in \mathbb{N}$, $a \otimes b$ is $a + b$
- for all diagrams $f : m_1 \rightarrow n_1, g : m_2 \rightarrow n_2$,
 $f \otimes g : m_1 \otimes m_2 \rightarrow n_1 \otimes n_2$ is juxtaposition



Definition (unit)

For all n , the arrow $\eta_n : 0 \rightarrow n + n$



is the *unit* (counit $\epsilon_n : n + n \rightarrow 0$ similar).

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Using this unit, all objects are self-dual.

Monoidal Functors

Definition

Let $\langle C_1, \otimes_1, 1_1, a_1, \iota_1 \rangle$ and $\langle C_2, \otimes_2, 1_2, a_2, \iota_2 \rangle$ be monoidal categories, $F : C_1 \rightarrow C_2$ be a functor with a natural isomorphism

$$J_{X,Y} : F(X) \otimes_2 F(Y) \xrightarrow{\sim} F(X \otimes_1 Y)$$

such that $F(1_1)$ is isomorphic to 1_2 and the diagram

$$\begin{array}{ccc} (F(X) \otimes^l F(Y)) \otimes^l F(Z) & \xrightarrow{a_{F(X), F(Y), F(Z)}^l} & F(X) \otimes^l (F(Y) \otimes^l F(Z)) \\ \downarrow J_{X,Y} \otimes^l \text{id}_{F(Z)} & & \downarrow \text{id}_{F(X)} \otimes^l J_{Y,Z} \\ F(X \otimes Y) \otimes^l F(Z) & & F(X) \otimes^l F(Y \otimes Z) \\ \downarrow J_{X \otimes Y, Z} & & \downarrow J_{X, Y \otimes Z} \\ F((X \otimes Y) \otimes Z) & \xrightarrow{F(a_{X,Y,Z})} & F(X \otimes (Y \otimes Z)) \end{array}$$

commutes, then $\langle F, H \rangle$ is a *monoidal functor*.

Example

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Forgetful functors are monoidal.